

Recent developments in finite alphabet rank \mathcal{S} -adic sequences

Rencontre IZES

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1. Definitions
2. Finite alphabet rank \mathcal{S} -adic subshifts
3. Recent developments
 - Factors
 - Asymptotic pairs
4. \mathcal{S} -adic representations

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- 4 We will consider only **minimal** subshifts X , *i.e.* s.t. every orbit $\text{orb}_S(x) := \{S^k x : k \in \mathbb{Z}\}$ is dense in X .

Definitions: \mathcal{S} -adic subshifts

① For $n \geq 0$, let $\tau_n: \mathcal{A}_{n+1} \rightarrow \mathcal{A}_n^+$ be a map, which we call **substitution**.

- We define the extension $\tau_n: \mathcal{A}_{n+1}^+ \rightarrow \mathcal{A}_n^+$ by concatenation:

$$\tau_n(a_1 a_2 \dots a_\ell) = \tau_n(a_1) \tau_n(a_2) \dots \tau_n(a_\ell)$$

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② Then, $\tau = (\tau_n)_{n \geq 0}$ is called an **\mathcal{S} -adic sequence** if

$$\lim_{n \rightarrow +\infty} \min_{a \in \mathcal{A}_n} |\tau_{[0,n]}(a)| = +\infty.$$

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③ We define the **n -th \mathcal{S} -adic subshift** $X_\tau^{(n)}$ as follows:

- $x \in X_\tau^{(n)}$ iff $\forall \ell \geq 0, \exists N > n, a \in \mathcal{A}_N$ s.t. $x_{[-\ell, \ell]}$ occurs in $\tau_{[n,N]}(a)$.

We set $X_\tau := X_\tau^{(0)}$.

Definitions: factorizations

- ① **Fact.** ¹ For any $n > 0$ and $x \in X_\tau$, there exists a $\tau_{[0,n]}$ -factorization of x in $X_\tau^{(n)}$, that is, $k \in \mathbb{Z}$ and $y \in X_\tau^{(n)}$ s.t.

$$x = S^k \tau_{[0,n]}(y) \text{ and } 0 \leq k < |\tau_{[0,n]}(y_0)|.$$

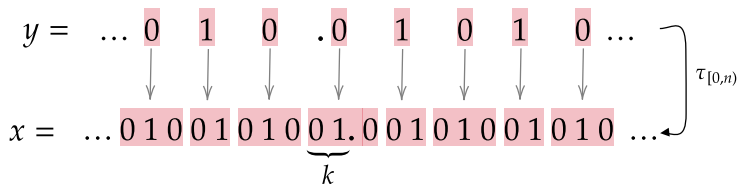


Figure: Here, $\tau_{[0,n]}: \{0, 1\} \rightarrow \{0, 1\}^+$ maps $0 \mapsto 010$ and $1 \mapsto 01$.

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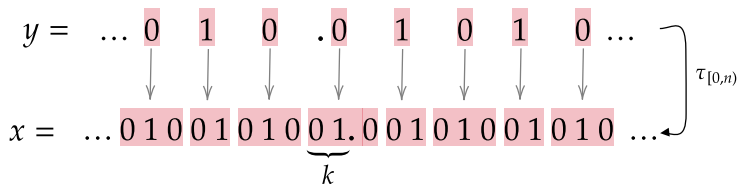


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- ② τ is **recognizable** if (k, y) is unique for all $n > 0$ and $x \in X_\tau$.

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Definitions: alphabet rank

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- 2 A **contraction** of τ is an \mathcal{S} -adic sequence of the form $\tau' = (\tau_{[0, n_1)}, \tau_{[n_1, n_2)}, \dots)$.
- 3 τ and τ' generate the same \mathcal{S} -adic subshift X_τ .

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- 4 If $\text{Alph-Rank}(\tau) < \infty$, then there is a contraction $\tau' = (\tau_{[n_k, n_{k+1})})_{k \geq 0}$ s.t. $\#\mathcal{A}_{n_k} = \text{Alph-Rank}(\tau)$, $\forall k \geq 1$.

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The topological rank

- ① **Framework:** minimal subshifts X_τ , where τ has finite alphabet rank.

²Corollary 1.4 in *Symbolic factors of S -adic subshifts of finite alphabet rank*, B. Espinoza, 2022.

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- 2 $\tau: \mathcal{A} \rightarrow \mathcal{B}^+$ is **proper** if $\exists a, b \in \mathcal{B}$ s.t. $\forall c \in \mathcal{A}$, $\tau(c)$ starts with a and ends with b .
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- 4 For X a minimal subshift, the following are equivalent²:
 - $X = X_\tau$ for some τ s.t. $\text{Alph-Rank}(\tau) < +\infty$.
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A minimal subshift X has **topological rank K** if K is the least number for which $X = X_\tau$ for some recognizable and proper τ s.t. $\text{Alph-Rank}(\tau) = K$.

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- 5 **Framework:** it is equivalent to work with finite top. rank subshifts.

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Examples of finite top. rank subshifts

- 1 substitutive and linearly recurrent subshifts;
- 2 Sturmian, (natural codings of) minimal IET, dendric subshifts;
- 3 some Toeplitz subshifts;
- 4 subshifts X whose word-complexity function p_X has linear or non-superlinear growth;
- 5 for any $f: \mathbb{N} \rightarrow \mathbb{N}$ s.t. $\frac{1}{n} \log f(n) \rightarrow 0$ there exists X with finite top. rank and $p_X(n)/f(n) \rightarrow +\infty$.

Known properties of finite top. rank subshifts

Although the finite top. rank class contains a broad spectrum of minimal subshifts, it is possible to prove general theorems for it.

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Although the finite top. rank class contains a broad spectrum of minimal subshifts, it is possible to prove general theorems for it. In a finite top.

rank subshift X_τ :

- 1 the topological entropy is zero (Handelman-Boyle, '94);
- 2 there are only finitely many ergodic measures, and they can be computed from τ (Bezuglyi *et. al.* '13);
- 3 the dimension group can be computed from τ (Herman *et. al.*, '99);
- 4 there is criteria stated in terms of τ for deciding if a given complex number is an eigenvalue (Durand *et. al.*, '19);
- 5 among other properties and tools.

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 - this is an equicontinuous system.
- 3 Example: (X, S) is a system.
 - this is an expansive system.
- 4 A **factor map** of X is a continuous map $\pi: X \rightarrow Y$ s.t. $T\pi(x) = \pi(Sx)$, $\forall x \in X$. Then, Y is a **factor** of X .
- 5 Factors are (topological) representations of X in Y .

Subshift factors

- 1 Suppose that Y is a subshift.
- 2 (Curtis–Hedlund–Lyndon theorem) There exists $r \geq 0$ and $\phi: \mathcal{A}^{2r+1} \rightarrow \mathcal{B}$ s.t. π is described by

$$x = (x_n)_{n \in \mathbb{Z}} \mapsto \pi(x) = (\phi(x_{[n-r, n+r]}))_{n \in \mathbb{Z}}.$$

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- 3 We first consider subshift factors of \mathcal{S} -adic subshifts.
- 4 **Our framework:** a minimal subshift $X = X_\tau$ s.t. $\tau = (\tau_n: \mathcal{A}_{n+1} \rightarrow \mathcal{A}_n^+)_{n \geq 0}$ is recognizable, proper, and of finite alphabet rank.

Subshift factors: \mathcal{S} -representation

Proposition (First representation of π)

Let $\pi: X_\tau \rightarrow Y$ be a subshift factor, with $Y \subseteq \mathcal{B}^{\mathbb{Z}}$. Then, there exist $\ell \geq 0$ and $\sigma: \mathcal{A}_\ell \rightarrow \mathcal{B}^+$ s.t.

- 1 If $x \in X_\tau$ and (k, y) is its $\tau_{[0, \ell)}$ -factorization in $X_\tau^{(\ell)}$, then $\pi(x) = S^k \sigma(y)$.
- 2 $|\sigma(a)| = |\tau_{[0, \ell)}(a)|$ for all $a \in \mathcal{A}_\ell$.

Conversely, if σ and ℓ satisfy (2) then (1) defines a factor map π .

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- 3 σ' is proper and $\text{Alph-Rank}(\sigma') = \text{Alph-Rank}(\tau)$.
- 4 Problem: σ' is recognizable iff π is a conjugacy.

Subshift factors: \mathcal{S} -representation

Theorem (E. 2021)

Up to a contraction of τ , there exist an \mathcal{S} -adic sequence

$\sigma = (\sigma_n: \mathcal{B}_{n+1} \rightarrow \mathcal{B}_n^+)_{n \geq 0}$ and morphisms $\phi_n: \mathcal{A}_n \rightarrow \mathcal{B}_n^+$ s.t.

- 1 $Y = X_\sigma$, σ is proper, **recognizable**, $\text{Alph-Rank}(\sigma) \leq \text{Alph-Rank}(\tau)$.
- 2 For $n \geq \ell$,

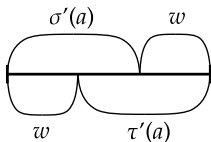
$$\begin{array}{ccc} X_\tau^{(\ell)} & \xrightarrow{\phi_\ell} & X_\sigma^{(\ell)} \\ \tau_{[0,\ell]} \downarrow & \searrow \sigma & \downarrow \sigma_{[0,\ell]} \\ X_\tau & \xrightarrow{\pi} & Y \end{array}$$

$$\begin{array}{ccc} X_\tau^{(n+1)} & \xrightarrow{\phi_{n+1}} & X_\sigma^{(n+1)} \\ \tau_n \downarrow & & \downarrow \sigma_n \\ X_\tau^{(n)} & \xrightarrow{\phi_n} & X_\sigma^{(n)} \end{array}$$

- 1 Corollary: the top. rank of Y is at most the one of X_τ .
- 2 The proof of the theorem is constructive and allows to compute σ and ϕ_n .

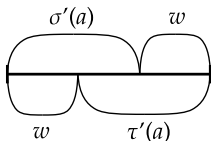
Subshift factors: application to substitutions

- ① $\sigma', \tau': \mathcal{A} \rightarrow \mathcal{B}^+$ are **rotationally conjugate** if there exists a word w s.t. $\sigma'(a)w = w\tau'(a)$ for all $a \in \mathcal{A}$.



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Proposition

Let $\tau: \mathcal{A} \rightarrow \mathcal{A}^+$, $\sigma: \mathcal{B} \rightarrow \mathcal{B}^+$ be primitive and proper substitutions. There exists a computable constant $n = n(\tau, \sigma)$ s.t. X_σ is a factor of X_τ iff there exist $\phi: \mathcal{A} \rightarrow \mathcal{B}^+$ and $p, q, r \leq n$ s.t.

- ① $\sigma^p \circ \phi$ and $\phi \circ \tau^q$ are rotationally conjugate;
- ② $|\tau^n(a)| = |\sigma^r \phi(a)|$ for all $a \in \mathcal{A}$.

Subshift factors: substitutive case

- 1 (Durand-Leroy '18) There is an algorithm that decides whether X_σ is a factor of X_τ .
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Question

Use this approach to continue exploring subshift factors of substitutive subshifts or other classes with finite top. rank.

More on subshift factors

Suppose that the minimal subshift X has finite topological rank.

- 1 Any self factor $\pi: X \rightarrow X$ is invertible (i.e. X is coalescent).
- 2 If $\pi: X \rightarrow Y$ is an infinite subshift factor, then $\exists k \in \mathbb{N}$ s.t. $\#\pi^{-1}(y) = k$ for almost all $y \in Y$.

Theorem

The number of subshift factors of X up to conjugacy is finite.

Other types of factors

Suppose that X is a minimal subshift with finite top. rank.

- ① **Proposition:** if $\pi: X \rightarrow Y$ is a factor and Y is totally disconnected, then Y is either a subshift or an odometer.

Other types of factors

Suppose that X is a minimal subshift with finite top. rank.

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Other types of factors

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Question

What can be said of factors $\pi: X \rightarrow Y$ s.t. Y is a **connected** topological space?

- This case correspond to a geometrical representation of X .
- It is opposite to the totally disconnected case.
- It is a direction that has not been explored even in the substitutive case.

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Asymptotic pairs

- 1 A (centered) **asymptotic pair** in X is a tuple $(x, y) \in X \times X$ s.t. $x_{(-\infty, 0)} = y_{(-\infty, 0)}$ and $x_0 \neq y_0$.

Theorem (E., Maass, 2020)

Let $\tau = (\tau_n: \mathcal{A}_{n+1} \rightarrow \mathcal{A}_n^+)_{n \geq 0}$ be such that

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Then, X_τ has finitely many asymptotic pairs.

- 2 This extends a previous result by Donoso *et. al.* from 2016.
- 3 The theorem can be used to bound the number of *automorphisms* and *symbolic factors* of X_τ in the minimal case.

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Other types of factors

Let X be a minimal subshift.

- 1 The **complexity function** $p_X: \mathbb{N} \rightarrow \mathbb{N}$ is defined by

$$p_X(n) = \mathbb{N}^\circ \text{ of words of length } n \text{ occurring in at least one } x \in X.$$

- 2 X has **linear growth** if

$$\limsup_{n \rightarrow +\infty} p_X(n)/n < +\infty.$$

- 3 X has **non-superlinear growth** if

$$\liminf_{n \rightarrow +\infty} p_X(n)/n < +\infty.$$

Linear growth subshifts

- The **root** $\text{root}(w)$ of a word w is the shortest word v such that $w = v^k$ for some $k \geq 1$.
- Example: $\text{root}(ababab) = ab$ and $\text{root}(ababababa) = ababababa$.

Theorem

A minimal subshift X has linear growth if and only if there exist $C \geq 1$ and an S -adic sequence $\tau = (\tau_n: \mathcal{A}_{n+1}^+ \rightarrow \mathcal{A}_n)_{n \geq 0}$ generating X such that for every $n \geq 1$:

- 1 $\text{root}(\tau_{[0,n]}(\mathcal{A}_n)) := \{\text{root}(\tau_{[0,n]}(a)) : a \in \mathcal{A}_n\}$ has at most C elements;
- 2 $|\tau_{[0,n]}(a)| \leq C \cdot |\tau_{[0,n]}(b)|$ for every $a, b \in \mathcal{A}_n$;
- 3 $|\tau_n(a)| \leq C$ for every $a \in \mathcal{A}_n$.

Linear growth subshifts

- 1 The previous characterization allows to recover known results about linear growth subshifts.
- 2 Cassaigne's Theorem: $p_X(n+1) - p_X(n)$ is bounded.
- 3 X has finitely many ergodic measures.
- 4 X has finite topological rank.

Non-superlinear growth subshifts

Theorem

A minimal subshift X has non-superlinear growth if and only if there exist $C \geq 1$ and an \mathcal{S} -adic sequence $\tau = (\tau_n: \mathcal{A}_{n+1}^+ \rightarrow \mathcal{A}_n)_{n \geq 0}$ generating X such that for every $n \geq 1$:

- 1 $\text{root}(\tau_{[0,n]}(\mathcal{A}_n)) = \{\text{root}(\tau_{[0,n]}(a)) : a \in \mathcal{A}_n\}$ has at most C elements;
- 2 $|\tau_{[0,n]}(a)| \leq C \cdot |\tau_{[0,n]}(b)|$ for every $a, b \in \mathcal{A}_n$.

Observations

- 1 The conditions depend on the unbounded product $\tau_{[0,n]}$, and thus are “non-local”.

Question

Is it possible to obtain a “local” characterization of subshifts with linear growth?

- 2 “Local” characterizations have been obtained for other classes, such as Sturmian, IET, and dendric.

Recent developments in finite alphabet rank \mathcal{S} -adic sequences

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